

# Part III — Extremal Graph Theory

## Theorems with proof

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Michaelmas 2017

These notes are not endorsed by the lecturers, and I have modified them (often significantly) after lectures. They are nowhere near accurate representations of what was actually lectured, and in particular, all errors are almost surely mine.

Turán's theorem, giving the maximum size of a graph that contains no complete  $r$ -vertex subgraph, is an example of an extremal graph theorem. Extremal graph theory is an umbrella title for the study of how graph and hypergraph properties depend on the values of parameters. This course builds on the material introduced in the Part II Graph Theory course, which includes Turán's theorem and also the Erdős–Stone theorem.

The first few lectures will cover the Erdős–Stone theorem and stability. Then we shall treat Szemerédi's Regularity Lemma, with some applications, such as to hereditary properties. Subsequent material, depending on available time, might include: hypergraph extensions, the flag algebra method of Razborov, graph containers and applications.

### **Pre-requisites**

A knowledge of the basic concepts, techniques and results of graph theory, as afforded by the Part II Graph Theory course, will be assumed. This includes Turán's theorem, Ramsey's theorem, Hall's theorem and so on, together with applications of elementary probability.

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## 1 The Erdős–Stone theorem

**Theorem** (Turán’s theorem). If  $G$  is a graph with  $|G| = n$ ,  $e(G) \geq t_r(n)$  and  $G \not\supseteq K_{r+1}$ . Then  $G = T_r(n)$ .

**Theorem** (Erdős–Stone, 1946). Let  $r \geq 1$  be an integer and  $\varepsilon > 0$ . Then there exists  $d = d(r, \varepsilon)$  and  $n_0 = n_0(r, \varepsilon)$  such that if  $|G| = n \geq n_0$  and

$$e(G) \geq \left(1 - \frac{1}{r} + \varepsilon\right) \binom{n}{2},$$

then  $G \supseteq K_{r+1}(t)$ , where  $t = \lfloor d \log n \rfloor$ .

**Lemma.** Let  $c, \varepsilon > 0$ . Then there exists  $n_1 = n_1(c, \varepsilon)$  such that if  $|G| = n \geq n_1$  and  $e(G) \geq (c + \varepsilon) \binom{n}{2}$ , then  $G$  has a subgraph  $H$  where  $\delta(H) \geq c|H|$  and  $|H| \geq \sqrt{\varepsilon n}$ .

*Proof.* The idea is that we can keep removing the vertex of smallest degree and then we must eventually get the  $H$  we want. Suppose this doesn’t give us a suitable graph even after we’ve got  $\sqrt{\varepsilon n}$  vertices left. That means we can find a sequence

$$G = G_n \supseteq G_{n-1} \supseteq G_{n-2} \supseteq \cdots G_s,$$

where  $s = \lfloor \varepsilon^{1/2} n \rfloor$ ,  $|G_j| = j$  and the vertex in  $G_j \setminus G_{j-1}$  has degree  $< cj$  in  $G_j$ .

We can then calculate

$$\begin{aligned} e(G_s) &> (c + \varepsilon) \binom{n}{2} - c \sum_{j=s+1}^n j \\ &= (c + \varepsilon) \binom{n}{2} - c \left\{ \binom{n+1}{2} - \binom{s+1}{2} \right\} \\ &\sim \frac{\varepsilon n^2}{2} \end{aligned}$$

as  $n$  gets large (since  $c$  and  $s$  are fixed numbers). In particular, this is  $> \binom{s}{2}$ . But  $G_s$  only has  $s$  vertices, so this is impossible.  $\square$

**Lemma.** Let  $r \geq 1$  be an integer and  $\varepsilon > 0$ . Then there exists a  $d_1 = d_1(r, \varepsilon)$  and  $n_2 = n_2(r, \varepsilon)$  such that if  $|G| = n \geq n_2$  and

$$\delta(G) \geq \left(1 - \frac{1}{r} + \varepsilon\right) n,$$

then  $G \supseteq K_{r+1}(t)$ , where  $t = \lfloor d_1 \log n \rfloor$ .

*Proof of Erdős–Stone theorem.* Provided  $n_0$  is large, say  $n_0 > n_1 \left(1 - \frac{1}{r} + \frac{\varepsilon}{2}, \frac{\varepsilon}{2}\right)$ , we can apply the first lemma to  $G$  to obtain a subgraph  $H \subseteq G$  where  $|H| > \sqrt{\frac{\varepsilon}{2}} n$ , and  $\delta(H) \geq \left(1 - \frac{1}{r} + \frac{\varepsilon}{2}\right) |H|$ .

We can then apply our final lemma as long as  $\sqrt{\frac{\varepsilon}{2}} n$  is big enough, and obtain  $K_{r+1}(t) \subseteq H \subseteq G$ , with  $t > \lfloor d_1(r, \varepsilon/2) \log(\sqrt{\frac{\varepsilon}{2}} n) \rfloor$ .  $\square$

*Proof of lemma.* We proceed by induction on  $r$ . If  $r = 0$  or  $\varepsilon \geq \frac{1}{r}$ , the theorem is trivial. Otherwise, by the induction hypothesis, we may assume  $G \supseteq K_r(T)$  for

$$T = \left\lfloor \frac{2t}{\varepsilon r} \right\rfloor.$$

Call it  $K = K_r(T)$ . This is always possible, as long as

$$d_1(r, \varepsilon) < \frac{\varepsilon r}{2} d_1\left(r-1, \frac{1}{r(r-1)}\right).$$

The exact form is not important. The crucial part is that  $\frac{1}{r(r-1)} = \frac{1}{r-1} - \frac{1}{r}$ , which is how we chose the  $\varepsilon$  to put into  $d_1(r-1, \varepsilon)$ .

Let  $U$  be the set of vertices in  $G - K$  having at least  $(1 - \frac{1}{r} + \frac{\varepsilon}{2})|K|$  neighbours in  $K$ . Calculating

$$\left(1 - \frac{1}{r} + \frac{\varepsilon}{2}\right)|K| = \left(1 - \frac{1}{r} + \frac{\varepsilon}{2}\right)rT = (r-1)T + \frac{\varepsilon r}{2}T \geq (r-1)T + t,$$

we see that every element in  $U$  is joined to at least  $t$  vertices in each class, and hence is joined to some  $K_r(t) \subseteq K$ .

So we have to show that  $|U|$  is large. If so, then by a pigeonhole principle argument, it follows that we can always find enough vertices in  $U$  so that adding them to  $K$  gives a  $K_{r+1}(t)$ .

**Claim.** There is some  $c > 0$  (depending on  $r$  and  $\varepsilon$ ) such that for  $n$  sufficiently large, we have

$$|U| \geq cn.$$

The argument to establish these kinds of bounds is standard, and will be used repeatedly. We write  $e(K, G - K)$  for the number of edges between  $K$  and  $G - K$ . By the minimum degree condition, each vertex of  $K$  sends at least  $(1 - \frac{1}{r} + \varepsilon)n - |K|$  edges to  $G - K$ . Then we have two inequalities

$$\begin{aligned} |K| \left\{ \left(1 - \frac{1}{r} + \varepsilon\right)n - |K| \right\} &\leq e(K, G - K) \\ &\leq |U||K| + (n - |U|) \left(1 - \frac{1}{r} + \frac{\varepsilon}{2}\right)|K|. \end{aligned}$$

Rearranging, this tells us

$$\frac{\varepsilon n}{2} - |K| \leq |U| \left( \frac{1}{r} - \frac{\varepsilon}{2} \right).$$

Now note that  $|K| \sim \log n$ , so for large  $n$ , the second term on the left is negligible, and it follows that  $U \sim n$ .

We now want to calculate the number of  $K_r(t)$ 's in  $K$ . To do so, we use the simple inequality

$$\binom{n}{k} \leq \left(\frac{en}{k}\right)^k.$$

Then we have

$$\#K_r(t) = \binom{T}{t}^r \leq \left(\frac{eT}{t}\right)^{rt} \leq \left(\frac{3e}{\varepsilon r}\right)^{rt} \leq \left(\frac{3e}{\varepsilon r}\right)^{rd \log n} = n^{rd \log(3e/\varepsilon r)}$$

Now if we pick  $d$  sufficiently small, then this tells us  $\#K_r(t)$  grows sublinearly with  $n$ . Since  $|U|$  grows linearly, and  $t$  grows logarithmically, it follows that for  $n$  large enough, we have

$$|U| \geq t \cdot \#K_r(t).$$

Then by the pigeonhole principle, there must be a set  $W \subseteq U$  of size  $t$  joined to the same  $K_r(t)$ , giving a  $K_{r+1}(t)$ .  $\square$

**Theorem** (Erdős–Simonovits). Let  $F$  be a fixed graph with chromatic number  $\chi(F) = r + 1$ . Then

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, F)}{\binom{n}{2}} = 1 - \frac{1}{r}.$$

*Proof.* Since  $\chi(F) = r + 1$ , we know  $F$  cannot be embedded in an  $r$ -partite graph. So in particular,  $F \not\subseteq T_r(n)$ . So

$$\text{ex}(n, F) \geq t_r(n) \geq \left(1 - \frac{1}{r}\right) \binom{n}{2}.$$

On the other hand, given any  $\varepsilon > 0$ , if  $|G| = n$  and

$$e(G) \geq \left(1 - \frac{1}{r} + \varepsilon\right) \binom{n}{2},$$

then by the Erdős–Stone theorem, we have  $G \supseteq K_{r+1}(|F|) \supseteq F$  provided  $n$  is large. So we know that for every  $\varepsilon > 0$ , we have

$$\limsup \frac{\text{ex}(n, F)}{\binom{n}{2}} \leq 1 - \frac{1}{r} + \varepsilon.$$

So we are done.  $\square$

**Theorem.** Given  $r \in \mathbb{N}$ , there exists  $\varepsilon_r > 0$  such that if  $0 < \varepsilon < \varepsilon_r$ , then there exists  $n_3(r, \varepsilon)$  so that if  $n > n_3$ , there exists a graph  $G$  of order  $n$  such that

$$e(G) \geq \left(1 - \frac{1}{r} + \varepsilon\right) \binom{n}{2}$$

but  $K_{r+1}(t) \not\subseteq G$ , where

$$t = \left\lceil \frac{3 \log n}{\log 1/\varepsilon} \right\rceil.$$

*Proof.* Start with the Turán graph  $T_r(n)$ . Let  $m = \lceil \frac{n}{r} \rceil$ . Then there is a class  $W$  of  $T_r(n)$  of size  $m$ . The strategy is to add  $\varepsilon \binom{n}{2}$  edges inside  $W$  to obtain  $G$  such that  $G[W]$  (the subgraph of  $G$  formed by  $W$ ) does not contain a  $K_2(t)$ . It then follows that  $G \not\supseteq K_{r+1}(t)$ , but also

$$e(G) \geq t_r(n) + \varepsilon \binom{n}{2} \geq \left(1 - \frac{1}{r} + \varepsilon\right) \binom{n}{2},$$

as desired.

To see that such an addition is possible, we choose edges inside  $W$  independently with probability  $p$ , to be determined later. Let  $X$  be the number of edges chosen, and  $Y$  be the number of  $K_2(t)$  created. If  $\mathbb{E}[X - Y] > \varepsilon \binom{n}{2}$ , then this

means there is an actual choice of edges with  $X - Y > \varepsilon \binom{n}{2}$ . We then remove an edge from each  $K_2(t)$  to leave a  $K_2(t)$ -free graph with at least  $X - Y > \varepsilon \binom{n}{2}$  edges.

So we want to actually compute  $\mathbb{E}[X - Y]$ . Seeing that asymptotically,  $m$  is much larger than  $t$ , we have

$$\begin{aligned} \mathbb{E}[X - Y] &= \mathbb{E}[X] - \mathbb{E}[Y] \\ &= p \binom{m}{2} - \frac{1}{2} \binom{m}{t} \binom{m-t}{t} p^{t^2} \\ &\sim \frac{1}{2} p m^2 - \frac{1}{2} m^{2t} p^{t^2} \\ &= \frac{1}{2} p m^2 (1 - m^{2t-2} p^{t^2-1}) \\ &= \frac{1}{2} p m^2 (1 - (m^2 p^{t+1})^{t-1}). \end{aligned}$$

We pick  $p = 3\varepsilon r^2$  and  $\varepsilon_r = (3r^2)^{-6}$ . Then  $p < \varepsilon^{5/6}$ , and

$$m^2 p^{t+1} < m^2 \varepsilon^{5/6(t+1)} < (m^2 m^{-5/2})^{t-1} < \frac{1}{2}.$$

Hence, we find that

$$\mathbb{E}[x - y] \geq \frac{1}{4} p m^2 > \varepsilon \binom{n}{2}. \quad \square$$

## 2 Stability

**Theorem.** Let  $t, r \geq 2$  be fixed, and suppose  $|G| = n$  and  $G \not\supseteq K_{r+1}(t)$ . If

$$e(G) = \left(1 - \frac{1}{r} + o(1)\right) \binom{n}{2}.$$

Then

- (i) There exists  $T_r(n)$  on  $V(G)$  with  $|E(G) \Delta E(T_r(n))| = o(n^2)$ .
- (ii)  $G$  contains an  $r$ -partite subgraph with  $(1 - \frac{1}{r} + o(1)) \binom{n}{2}$  edges.
- (iii)  $G$  contains an  $r$ -partite subgraph with minimum degree  $(1 - \frac{1}{r} + o(1))n$ .

*Proof.* We first prove (ii), which looks the weakest. We will then argue that the others follow from this (and in fact they are equivalent).

By Erdős–Stone, we can always find some  $K_r(s) = K \subseteq G$  for some  $s = s(n) \rightarrow \infty$ . We can, of course, always choose  $s(n)$  so that  $s(n) \leq \log n$ , which will be helpful for our future analysis. We let  $C_i$  be the  $i$ th class of  $K$ .

- By throwing away  $o(n)$  vertices, we may assume that

$$\delta(G) \geq \left(1 - \frac{1}{r} + o(1)\right)n$$

- Let  $X$  be the set of all vertices joined to at least  $t$  vertices in each  $C_i$ . Note that there are  $\binom{s}{t}^r$  many copies of  $K_r(t)$ 's in  $K$ . So by the pigeonhole principle, we must have  $|X| < t \binom{s}{t}^r = o(n)$ , or else we can construct a  $K_{r+1}(t)$  inside  $G$ .
- Let  $Y$  be the set of vertices joined to fewer than  $(r-1)s - \frac{s}{2t} + t$  other vertices. By our bound on  $\delta(G)$ , we certainly have

$$e(K, G - K) \geq s(r-1 + o(1))n.$$

On the other hand, every element of  $G$  is certainly joined to at most  $(r-1)s + t$  edges, since  $X$  was thrown away. So we have

$$((r-1)s + t)(n - |Y|) + |Y| \left( (r-1)s - \frac{s}{2t} + t \right) \geq e(K, G - K).$$

So we deduce that  $|Y| = o(n)$ . Throw  $Y$  away.

Let  $V_i$  be the set of vertices in  $G \setminus K$  joined to fewer than  $t$  of  $C_i$ . It is now enough to show that  $e(G[V_i]) = o(n^2)$ . We can then throw away the edges in each  $V_i$  to obtain an  $r$ -partite graph.

Suppose on the contrary that  $e(G[V_j]) \geq \varepsilon n^2$ , say, for some  $j$ . Then Erdős–Stone with  $r = 1$  says we have  $G[V_j] \supseteq K_2(t)$ . Each vertex of  $K_2(t)$  has at least  $s - \frac{s}{2t} + 1$  in each other  $C_i$ , for  $i \neq j$ . So we see that  $K_2(t)$  has at least  $s - 2t \left(\frac{s}{2t} - 1\right) > t$  common neighbours in each  $C_i$ , giving  $K_{r+1}(t) \subseteq G$ .

It now remains to show that (i) and (iii) follows from (ii). The details are left for the example sheet, but we sketch out the rough argument. (ii)  $\Rightarrow$  (i) is

straightforward, since an  $r$ -partite graph with that many edges is already pretty close to being a Turán graph.

To deduce (iii), since we can assume  $\delta(G) \geq (1 - \frac{1}{r} + o(1))n$ , we note that if (iii) did not hold, then we have  $\varepsilon n$  vertices of degree  $\leq (1 - \frac{1}{r} - \varepsilon)n$ , and their removal leaves a graph of order  $(1 - \varepsilon)n$  with at least

$$\left(1 - \frac{1}{r} + o(1)\right) \binom{n}{2} - \varepsilon n \left(1 - \frac{1}{r} - \varepsilon\right) n > \left(1 - \frac{1}{r} + \varepsilon^2\right) \binom{(1 - \varepsilon)n}{2},$$

which by Erdős–Stone would contain  $K_{r+1}(t)$ . So we are done.  $\square$

**Corollary.** Let  $\chi(F) = r + 1$ , and let  $G$  be extremal for  $F$ , i.e.  $G \not\supseteq F$  and  $e(G) = \text{ex}(|G|, F)$ . Then  $\delta(G) = (1 - \frac{1}{r} + o(1))n$ .

*Proof.* By our asymptotic bound on  $\text{ex}(n, F)$ , it cannot be that  $\delta(G)$  is greater than  $(1 - \frac{1}{r} + o(1))n$ . So there must be a vertex  $v \in G$  with  $d(v) \leq (1 - \frac{1}{r} - \varepsilon)n$ .

We now apply version (iii) of the stability theorem to obtain an  $r$ -partite subgraph  $H$  of  $G$ . We can certainly pick  $|F|$  vertices in the same part of  $H$ , which are joined to  $m = (1 - \frac{1}{r} + o(1))$  common neighbours. Form  $G^*$  from  $G - v$  by adding a new vertex  $u$  joined to these  $m$  vertices. Then  $e(G^*) > e(G)$ , and so the maximality of  $G$  entails  $G^* \supseteq F$ .

Pick a copy of  $F$  in  $G^*$ . This must involve  $u$ , since  $G$ , and hence  $G - v$  does not contain a copy of  $F$ . Thus, there must be some  $x$  amongst the  $|F|$  vertices mentioned. But then we can replace  $u$  with  $x$  and we are done.  $\square$

**Theorem (Simonovits).** Let  $F$  be  $(r + 1)$ -edge-critical, i.e.  $\chi(F) = r + 1$  but  $\chi(F \setminus e) = r$  for every edge  $e$  of  $F$ . Then for large  $n$ ,

$$\text{ex}(n, F) = t_r(n).$$

and the only extremal graph is  $T_r(n)$ .

*Proof.* Let  $G$  be an extremal graph for  $F$  and let  $H$  be an  $r$ -partite subgraph with

$$\delta(H) = \left(1 - \frac{1}{r} + o(1)\right)n.$$

Note that  $H$  necessarily have  $r$  parts of size  $(\frac{1}{r} + o(1))n$  each. Assign each of the  $o(n)$  vertices in  $V(G) \setminus V(H)$  to a class where it has fewest neighbours.

Suppose some vertex  $v$  has  $\varepsilon n$  neighbours in its own class. Then it has  $\geq \varepsilon n$  in each class. Pick  $\varepsilon n$  neighbours of  $v$  in each class of  $H$ . These neighbours span a graph with at least  $(1 - \frac{1}{r} + o(1)) \binom{r\varepsilon n}{2}$  edges. So by Erdős–Stone (or arguing directly), they span  $F - w$  for some vertex  $w$  (contained in  $K_r(|F|)$ ). Hence  $G \supseteq F$ , contradiction.

Thus, each vertex of  $G$  has only  $o(n)$  vertices in its own class. So it is joined to all but  $o(n)$  vertices in every other class.

Suppose some class of  $G$  contains an edge  $xy$ . Pick a set  $Z$  of  $|F|$  vertices in that class with  $\{x, y\} \in Z$ . Now  $Z$  has  $(\frac{1}{r} + o(1))n$  common neighbours in each class, so (by Erdős–Stone or directly) these common neighbours span a  $K_{r-1}(|F|)$ . But together with  $Z$ , we have a  $K_r(|F|)$  but with an edge added inside a class. But by our condition that  $F \setminus e$  is  $r$ -partite for any  $e$ , this subgraph contains an  $F$ . This is a contradiction.



So we conclude that no class of  $G$  contains an edge. So  $G$  itself is  $r$ -partite, but the  $r$ -partite graph with most edges is  $T_r(n)$ , which does not contain  $F$ . So  $G = T_r(n)$ .  $\square$

### 3 Supersaturation

**Theorem** (Erdős–Simonovits). Let  $H$  be some  $\ell$ -uniform hypergraph. Then for all  $\varepsilon > 0$ , there exists  $\delta(H, \varepsilon)$  such that every  $\ell$ -uniform hypergraph  $G$  with  $|G| = n$  and

$$e(G) > (\pi(H) + \varepsilon) \binom{n}{\ell}$$

contains  $\lfloor \delta n^{|H|} \rfloor$  copies of  $H$ .

*Proof.* For each  $m$ -set  $M \subseteq V(G)$ , we let  $G[M]$  be the sub-hypergraph induced by these vertices. Let the number of subsets  $M$  with  $e(G[M]) > (\pi(H) + \frac{\varepsilon}{2}) \binom{m}{\ell}$  be  $\eta \binom{n}{m}$ . Then we can estimate

$$\begin{aligned} (\pi(H) + \varepsilon) \binom{n}{\ell} &\leq e(G) = \frac{\sum_M e(G[M])}{\binom{n-\ell}{m-\ell}} \\ &\leq \frac{\eta \binom{n}{m} \binom{m}{\ell} + (1 - \eta) \binom{n}{m} (\pi(H) + \frac{\varepsilon}{2}) \binom{m}{\ell}}{\binom{n-\ell}{m-\ell}}. \end{aligned}$$

So if  $n > m$ , then

$$\pi(H) + \varepsilon \leq \eta + (1 + \eta) \left( \pi(H) + \frac{\varepsilon}{2} \right).$$

So

$$\eta \geq \frac{\frac{\varepsilon}{2}}{1 - \pi(H) - \frac{\varepsilon}{2}} > 0.$$

The point is that it is positive, and that's all we care about.

We pick  $m$  large enough so that

$$\text{ex}(m, H) < \left( \pi(H) + \frac{\varepsilon}{2} \right) \binom{m}{\ell}.$$

Then  $H \subseteq G[M]$  for at least  $\eta \binom{n}{m}$  choices of  $M$ . Hence  $G$  contains at least

$$\frac{\eta \binom{n}{m}}{\binom{n-|H|}{m-|H|}} = \frac{\eta \binom{n}{|H|}}{\binom{m}{|H|}}$$

copies of  $H$ , and we are done. (Our results hold when  $n$  is large enough, but we can choose  $\delta$  small enough so that  $\delta n^{|H|} < 1$  when  $n$  is small)  $\square$

**Theorem** (Lorden, 1961). Let  $G$  have degree sequence  $d_1, \dots, d_n$ . Then

$$k_3(G) + k_3(\bar{G}) = \binom{n}{3} - (n-2)e(G) + \sum_{i=1}^n \binom{d_i}{2}.$$

*Proof.* The number of paths of length 2 in  $G$  and  $\bar{G}$  is precisely

$$\sum_{i=1}^n \left( \binom{d_i}{2} + \binom{n-1-d_i}{2} \right) = 2 \sum_{i=1}^n \binom{d_i}{2} - 2(n-2)e(G) + 3 \binom{n}{3},$$

since to find a path of length 2, we pick the middle vertex and then pick the two edges. A complete or empty  $K_3$  contains 3 such paths; Other sets of three vertices contain exactly 1 such path. Hence

$$\binom{n}{3} + 2(k_3(G) + k_3(\bar{G})) = \text{number of paths of length 2.} \quad \square$$

**Corollary** (Goodman, 1959). We have

$$k_3(G) + k_3(\bar{G}) \geq \frac{1}{24}n(n-1)(n-5).$$

*Proof.* Let  $m = e(G)$ . Then

$$k_3(G) + k_3(\bar{G}) \geq \binom{n}{3} - (n-2)m + n \binom{2m/n}{2}.$$

Then minimize over  $m$ .  $\square$

**Corollary.** For  $m = e(G)$  and  $n = |G|$ , we have

$$k_3(G) \geq \frac{m}{3n}(4m - n^2).$$

*Proof.* By Lorden's theorem, we know

$$k_3(G) + k_3(\bar{G}) = \binom{n}{3} - (n-2)e(\bar{G}) + \sum \binom{\bar{d}_i}{2},$$

where  $\bar{d}_i$  is the degree sequence in  $\bar{G}$ . But

$$3k_3(\bar{G}) \leq \sum \binom{\bar{d}_i}{2},$$

since the sum is counting the number of paths of length 2. So we find that

$$k_3(G) \geq \binom{n}{3} - (n-2)\bar{m} - \frac{2}{3}n \binom{2\bar{m}/n}{2},$$

and  $\bar{m} = \binom{n}{2} - m$ .  $\square$

**Theorem.** Let  $G$  be a graph. For any graph  $F$ , let  $i_F(G)$  be the number of induced copies of  $F$  in  $G$ , i.e. the number of subsets  $M \subseteq V(G)$  such that  $G[M] \cong F$ . So, for example,  $i_{K_p}(G) = k_p(G)$ .

Define

$$f(G) = \sum_F \alpha_F i_F(G),$$

with the sum being over a finite collection of graphs  $F$ , each being complete multipartite, with  $\alpha_F \in \mathbb{R}$  and  $\alpha_F \geq 0$  if  $F$  is not complete. Then amongst graphs of given order,  $f(G)$  is maximized on a complete multi-partite graph. Moreover, if  $\alpha_{\bar{K}_3} > 0$ , then there are no other maxima.

*Proof.* We may suppose  $\alpha_{\bar{K}_3} > 0$ , because the case of  $\alpha_{\bar{K}_3} = 0$  follows from a limit argument. Choose a graph  $G$  maximizing  $f$  and suppose  $G$  is not complete multipartite. Then there exist non-adjacent vertices  $x, y$  whose neighbourhoods  $X, Y$  differ.

There are four contributions to  $i_F(G)$ , coming from

- (i)  $F$ 's that contain both  $x$  and  $y$ ;
- (ii)  $F$ 's that contain  $y$  but not  $x$ ;
- (iii)  $F$ 's that contain  $x$  but not  $y$ ;
- (iv)  $F$ 's that contain neither  $x$  nor  $y$ .

We may assume that the contribution from (iii)  $\geq$  (ii), and if they are equal, then  $|X| \leq |Y|$ .

Consider what happens when we remove all edges between  $y$  and  $Y$ , and add edges from  $y$  to everything in  $X$ . Clearly (iii) and (iv) are unaffected.

- If (iii)  $>$  (ii), then after this move, the contribution of (ii) becomes equal to the contribution of (iii) (which didn't change), and hence strictly increased.

The graphs  $F$  that contribute to (i) are not complete, so  $\alpha_F \geq 0$ . Moreover, since  $F$  is complete multi-partite, it cannot contain a vertex in  $X \Delta Y$ . So making the move can only increase the number of graphs contributing to (i), and each contribution is non-negative.

- If (iii) = (ii) and  $|X| \leq |Y|$ , then we make similar arguments. The contribution to (ii) is unchanged this time, and we know the contribution of (i) strictly increased, because the number of  $\bar{K}_3$ 's contributing to (i) is the number of points not connected to  $x$  and  $y$ .

In both cases, the total sum increased.  $\square$

**Theorem** (Bollobás, 1976). Let  $1 \leq p \leq r$ , and for  $0 \leq x \leq \binom{n}{p}$ , let  $\psi(x)$  be a maximal convex function lying below the points

$$\{(k_p(T_q(n)), k_r(T_q(n))) : q = r - 1, r, \dots\} \cup \{(0, 0)\}.$$

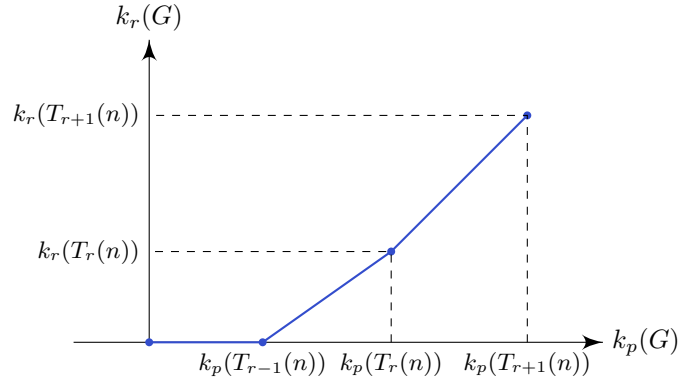
Let  $G$  be a graph of order  $n$ . Then

$$k_r(G) \geq \psi(k_p(G)).$$

*Proof.* Let  $f(G) = k_p(G) - ck_r(G)$ , where  $c > 0$ .

**Claim.** It is enough to show that  $f$  is maximized on a Turán graph for any  $c$ .

Indeed, suppose we plot out the values of  $(k_p(T_q(n)), k_r(T_q(n)))$ :



If the theorem doesn't hold, then we can pick a  $G$  such that  $(k_p(G), k_r(G))$  lies below  $\psi$ . Draw a straight line through the point with slope parallel to  $\psi$ . This has slope  $\frac{1}{c} > 0$  for some  $c$ . The intercept on the  $x$ -axis is then  $k_p(G) - ck_r(G)$ , which would be greater than  $f$  (any Turán graph) by convexity, a contradiction.

Now the previous theorem immediately tells us  $f$  is maximized on some complete multi-partite graph. Suppose this has  $q$  class, say of sizes  $a_1 \leq a_2 \leq \dots \leq a_q$ . It is easy to verify  $q \geq r - 1$ . In fact, we may assume  $q \geq r$ , else the maximum is on a Turán graph  $T_{r-1}(n)$ .

Then we can write

$$f(G) = a_1 a_q A - c a_q a_q B + C,$$

where  $A, B, C$  are rationals depending only on  $a_2, \dots, a_{q-1}$  and  $a_1 + a_q$  ( $A$  and  $B$  count the number of ways to pick a  $K_p$  and  $K_r$  respectively in a way that involves terms in both the first and last classes).

We wlog assume  $c$  is irrational. Hence  $a_q a_q A - c a_1 a_q = a_1 a_q (A - cB) \neq 0$ .

- If  $A - cB < 0$ , replace  $a_1$  and  $a_q$  by 0 and  $a_1 + a_q$ . This would then increase  $f$ , which is impossible.
- If  $A - cB > 0$  and  $a_1 \leq a_q - 2$ , then we can replace  $a_1, a_q$  by  $a_1 + 1, a_q - 1$  to increase  $f$ .

Hence  $a_1 \geq a_q - 1$ . So  $G = T_q(n)$ . □

## 4 Szemerédi's regularity lemma

**Lemma.** Let  $(U, W)$  be an  $\varepsilon$ -uniform pair with  $d(U, W) = d$ . Then

$$\begin{aligned} |\{u \in U : |\Gamma(u) \cap W| > (d - \varepsilon)|W|\}| &\geq (1 - \varepsilon)|U| \\ |\{u \in U : |\Gamma(u) \cap W| < (d + \varepsilon)|W|\}| &\geq (1 - \varepsilon)|U|, \end{aligned}$$

where  $\Gamma(u)$  is the set of neighbours of  $u$ .

*Proof.* Let

$$X = \{u \in U : |\Gamma(u) \cap W| \leq (d - \varepsilon)|W|\}.$$

Then  $e(X, W) \leq (d - \varepsilon)|X||W|$ . So

$$d(X, W) \leq d - \varepsilon = d(U, W) - \varepsilon.$$

So it fails the uniformity condition. Since  $W$  is definitely not small, we must have  $|X| < \varepsilon|U|$ .

The other case is similar, or observe that the complementary bipartite graph between  $U$  and  $W$  has density  $1 - d$  and is  $\varepsilon$ -uniform.  $\square$

**Lemma** (Graph building lemma). Let  $G$  be a graph containing distinct vertex subsets  $V_1, \dots, V_r$  with  $|V_i| = u$ , such that  $(V_i, V_j)$  is  $\varepsilon$ -uniform and  $d(V_i, V_j) \geq \lambda$  for all  $1 \leq i < j \leq r$ .

Let  $H$  be a graph with maximum degree  $\Delta(H) \leq \Delta$ . Suppose  $H$  has an  $r$ -colouring in which no colour is used more than  $s$  times, i.e.  $H \subseteq K_r(s)$ , and suppose  $(\Delta + 1)\varepsilon \leq \lambda^\Delta$  and  $s \leq \lfloor \varepsilon u \rfloor$ . Then  $H \subseteq G$ .

*Proof.* We wlog assume  $V(H) = \{1, \dots, k\}$ , and let  $c : V(H) \rightarrow \{1, \dots, r\}$  be a colouring of  $V(H)$  using no colour more than  $s$  times. We want to pick vertices  $x_1, \dots, x_k$  in  $G$  so that  $x_i x_j \in E(G)$  if  $ij \in E(H)$ .

We claim that, for  $0 \leq \ell \leq k$ , we can choose distinct vertices  $x_1, \dots, x_\ell$  so that  $x_j \in C_{c(j)}$ , and for  $\ell < j \leq k$ , a set  $X_j^\ell$  of candidates for  $x_j$  such that

- (i)  $X_j^\ell \subseteq V_{c(j)}$ ;
- (ii)  $x_i y_j \in E(G)$  for all  $y_j \in X_j^\ell$  and  $i \leq \ell$  such that  $ij \in E(H)$ .
- (iii)  $|X_j^\ell| \geq (\lambda - \varepsilon)^{|N(j, \ell)|} |V_{c(j)}|$ , where

$$N(j, \ell) = \{x_i : 1 \leq i \leq \ell \text{ and } ij \in E(H)\}.$$

The claim holds for  $\ell = 0$  by taking  $X_j^0 = V_{c(j)}$ .

By induction, suppose it holds for  $\ell$ . To pick  $x_{\ell+1}$ , of course we should pick it from our candidate set  $X_{\ell+1}^\ell$ . Then the first condition is automatically satisfied. Define the set

$$T = \{j > \ell + 1 : (\ell + 1)j \in E(H)\}.$$

Then each  $t \in T$  presents an obstruction to (ii) and (iii) being satisfied. To satisfy (ii), for each  $t \in T$ , we should set

$$X_t^{\ell+1} = X_t^\ell \cap \Gamma(x_{\ell+1}).$$

Thus, to satisfy (iii), we want to exclude those  $x_{\ell+1}$  that make this set too small. We define

$$Y_t = \left\{ y \in X_{\ell+1}^\ell : |\Gamma(y) \cap X_t^\ell| \leq (\lambda - \varepsilon)|X_t^\ell| \right\}.$$

So we want to find something in  $X_{\ell+1}^\ell \setminus \bigcup_{t \in T} Y_t$ . We also cannot choose one of the  $x_i$  already used. So our goal is to show that

$$\left| X_{\ell+1}^\ell - \bigcup_{t \in T} Y_t \right| > s - 1.$$

This will follow simply from counting the sizes of  $|X_{\ell+1}^\ell|$  and  $|Y_t|$ . We already have a bound on the size of  $|X_{\ell+1}^\ell|$ , and we shall show that if  $|Y_t|$  is too large, then it violates  $\varepsilon$ -uniformity.

Indeed, by definition of  $Y_t$ , we have

$$d(Y_t, X_{\ell+1}^\ell) \leq \lambda - \varepsilon \leq d(V_{c(t)}, V_{c(\ell+1)}) - \varepsilon.$$

So either  $|X_{\ell+1}^\ell| < \varepsilon|V_{c(t)}|$  or  $|Y_t| < \varepsilon|V_{c(\ell+1)}|$ . But the first cannot occur.

Indeed, write  $m = |N(\ell + 1, \ell)|$ . Then  $m + |T| \leq \Delta$ . In particular, since the  $|T| = 0$  case is trivial, we may assume  $m \leq \Delta - 1$ . So we can easily bound

$$|X_{\ell+1}^\ell| \geq (\lambda - \varepsilon)^{\Delta-1} |V_{c(t)}| \geq (\lambda^{\Delta-1} - (\Delta - 1)\varepsilon) |V_{c(t)}| > \varepsilon |V_{c(t)}|.$$

Thus, by  $\varepsilon$ -uniformity, it must be the case that

$$|Y_t| \leq \varepsilon |V_{c(\ell+1)}|.$$

Therefore, we can bound

$$\begin{aligned} \left| X_{\ell+1}^\ell - \bigcup_{t \in T} Y_t \right| &\geq (\lambda - \varepsilon)^m |V_{c(\ell+1)}| - (\Delta - m)\varepsilon |V_{c(\ell+1)}| \\ &\geq (\lambda^m - m\varepsilon - (\Delta - m)\varepsilon)u \geq \varepsilon u > s - 1. \end{aligned}$$

So we are done. □

**Corollary.** Let  $H$  be a graph with vertex set  $\{v_1, \dots, v_r\}$ . Let  $0 < \lambda, \varepsilon < 1$  satisfy  $r\varepsilon \leq \lambda^{r-1}$ .

Let  $G$  be a graph with disjoint vertex subsets  $V_1, \dots, V_r$ , each of size  $u \geq 1$ . Suppose each pair  $(V_i, V_j)$  is  $\varepsilon$  uniform, and  $d(V_i, V_j) \geq \lambda$  if  $v_i v_j \in E(H)$ , and  $d(V_i, V_j) \leq 1 - \lambda$  if  $v_i v_j \notin E(H)$ . Then there exists  $x_i \in V_i$  so that the map  $v_i \rightarrow x_i$  is an isomorphism  $H \rightarrow G[\{x_1, \dots, x_r\}]$ .

*Proof.* By replacing the  $V_i$ - $V_j$  edges by the complementary set whenever  $v_i v_j \notin E(H)$ , we may assume  $d(V_i, V_j) \geq \lambda$  for all  $i, j$ , and  $H$  is a complete graph.

We then apply the previous lemma with  $\Delta = r - 1$  and  $s = 1$ . □

**Theorem** (Szemerédi's regularity lemma). Let  $0 < \varepsilon < 1$  and let  $\ell$  be some natural number. Then there exists some  $L = L(\ell, \varepsilon)$  such that every graph has an  $\varepsilon$ -uniform equipartition into  $m$  parts for some  $\ell \leq m \leq L$ , depending on the graph.

**Lemma.** Let  $U' \subseteq U$  and  $W' \subseteq W$ , where  $|U'| \geq (1 - \delta)|U|$  and  $|W'| \geq (1 - \delta)|W|$ . Then

$$|d(U', W') - d(U, W)| \leq 2\delta.$$

*Proof.* Let  $d = d(U, W)$  and  $d' = d(U', W')$ . Then

$$d = \frac{e(U, W)}{|U||W|} \geq \frac{e(U', W')}{|U||W|} = d' \frac{|U'||W'|}{|U||W|} \geq d'(1 - \delta)^2.$$

Thus,

$$d' - d \leq d'(1 - (1 - \delta)^2) \leq 2\delta d' \leq 2\delta.$$

The other inequality follows from considering the complementary graph, which tells us

$$(1 - d') - (1 - d) \leq 2\delta. \quad \square$$

**Lemma.** Let  $x_1, \dots, x_n$  be real numbers with

$$X = \frac{1}{n} \sum_{i=1}^n x_i,$$

and let

$$x = \frac{1}{m} \sum_{i=1}^m x_i.$$

Then

$$\frac{1}{n} \sum_{i=1}^n x_i^2 \geq X^2 + \frac{m}{n-m} (x - X)^2 \geq X^2 + \frac{m}{n} (x - X)^2.$$

*Proof.* We have

$$\begin{aligned} \frac{1}{n} \sum_{i=1}^n x_i^2 &= \frac{1}{n} \sum_{i=1}^m x_i^2 + \frac{1}{n} \sum_{i=m+1}^n x_i^2 \\ &\geq \frac{m}{n} x^2 + \frac{n-m}{n} \left( \frac{nX - mx}{n-m} \right)^2 \\ &\geq X^2 + \frac{m}{n-m} (x - X)^2 \end{aligned}$$

by two applications of Cauchy–Schwarz. □

*Proof.* Define the index  $\text{ind}(\mathcal{P})$  of an equipartition  $\mathcal{P}$  into  $k$  parts  $V_i$  to be

$$\text{ind}(\mathcal{P}) = \frac{1}{k^2} \sum_{i < j} d^2(V_i, V_j).$$

We show that if  $\mathcal{P}$  is not  $\varepsilon$ -uniform, then there is a refinement equipartition  $\mathcal{Q}$  into  $k4^k$  parts, with  $\text{ind}(\mathcal{Q}) \geq \text{ind}(\mathcal{P}) + \frac{\varepsilon^5}{8}$ .

This is enough to prove the theorem. For choose  $t \geq \ell$  with  $4^t \varepsilon^5 \geq 100$ . Define recursively a function  $f$  by

$$f(0) = t, \quad f(j+1) = f(j)4^{f(j)}.$$



Let

$$N = f(\lceil 4\varepsilon^{-5} \rceil),$$

and pick  $L = N16^N$ .

Then, if  $n \leq L$ , then just take an equipartition into single vertices. Otherwise, begin with a partition into  $t$  parts. As long as the current partition into  $k$  parts is not  $\varepsilon$  uniform, replace it a refinement into  $4k^4$  parts. The point is that  $\text{ind}(\mathcal{P}) \leq \frac{1}{2}$  for any partition. So we can't do this more than  $4\varepsilon^{-5}$  times, at which point we have partitioned into  $N \leq L$  parts.

Note that the reason we had to set  $L = N16^N$  is that in our proof, we want to assume we have many vertices lying around.

The proof really is just one line, but students tend to complain about such short proofs, so let's try to explain it in a bit more detail. If the partition is not  $\varepsilon$ -uniform, this means we can further partition each part into uneven pieces. Then our previous lemma tells us this discrepancy allows us to push up  $\frac{1}{n} \sum x_i^2$ .

So given an equipartition  $\mathcal{P}$  that is not  $\varepsilon$ -uniform, for each non-uniform pair  $(V_i, V_j)$  of  $\mathcal{P}$ , we pick witness sets

$$X_{ij} \subseteq V_i, \quad X_{ji} \subseteq V_j$$

with  $|X_{ij}| \geq \varepsilon|V_i|$ ,  $|X_{ji}| \geq |V_j|$  and  $|d(X_{ij}, X_{ji}) - D(V_i, V_j)| \geq \varepsilon$ .

Fix  $i$ . Then the sets  $X_{ij}$  partition  $V_i$  into at most  $2^{k-1}$  atoms. Let  $m = \lfloor \frac{n}{k4^k} \rfloor$ , and let  $n = k4^k m + ak + b$ , where  $0 \leq a \leq 4^k$  and  $b \leq k$ . Then we see that

$$\lfloor n/k \rfloor = 4^k m + a$$

and the parts of  $\mathcal{P}$  have size  $4^k m + a$  or  $4^k m + a + 1$ , with  $b$  of the larger size.

Partition each part of  $\mathcal{P}$  into  $4^k$  sets, of size  $m$  or  $m + 1$ . The smaller  $V_i$  having  $a$  parts of size  $m + 1$ , and the larger having  $a + 1$  such pairs.

We see that any such partition is an equipartition into  $k4^k$  parts of size  $m$  or  $m + 1$ , with  $ak + b$  parts of larger size  $m + 1$ .

Let's choose such an equipartition  $\mathcal{Q}$  with parts as nearly as possible inside atoms, so each atom is a union of parts of  $\mathcal{Q}$  with at most  $m$  extra vertices.

All that remains is to check that  $\text{ind}(\mathcal{Q}) \geq \text{ind}(\mathcal{P}) + \frac{\varepsilon^5}{8}$ .

Let the sets of  $\mathcal{Q}$  within  $V_i$  be  $V_i(s)$ , where  $1 \leq s \leq 4^k \equiv q$ . So

$$V_i = \bigcup_{s=1}^q V_i(s).$$

Now

$$\sum_{1 \leq s, t \leq q} e(V_i(s), V_j(t)) = e(V_i, V_j).$$

We'd like to divide by some numbers and convert these to densities, but this is where we have to watch out. But this is still quite easy to handle. We have

$$\frac{m}{m+1} \leq q|V_i(s)| \leq |V_i| \leq \frac{m+1}{m} q|V_i(s)|$$

for all  $s$ . So we want  $m$  to be large for this to not hurt us too much.

So

$$\left( \frac{m}{m+1} \right) d(V_i, V_j) \leq \frac{1}{q^2} \sum_{s,t} d(V_i(s), V_j(t)) \leq \left( \frac{m+1}{m} \right)^2 d(V_i, V_j).$$

Using  $n \geq k16^k$ , and hence

$$\left(\frac{m}{m+1}\right)^2 \geq 1 - \frac{2}{m} \geq 1 - \frac{2}{4^k} \geq 1 - \frac{\varepsilon^5}{50},$$

we have

$$\left| \frac{1}{q^2} \sum_{s,t} d(V_i(s), V_j(t)) - d(V_i, V_j) \right| \leq \frac{\varepsilon^5}{49} + d(V_i, V_j).$$

In particular,

$$\frac{1}{q^2} \sum d^2(V_i(s), V_j(t)) \geq d^2(V_i, V_j) - \frac{\varepsilon^5}{25}.$$

The lower bound can be improved if  $(V_i, V_j)$  is not  $\varepsilon$ -uniform.

Let  $X_{ij}^*$  be the largest subset of  $X_{ij}$  that is the union of parts of  $\mathcal{Q}$ . We may assume

$$X_{ij}^* = \bigcup_{1 \leq s \leq r_i} V_i(s).$$

By an argument similar to the above, we have

$$\frac{1}{r_i r_j} \sum_{\substack{1 \leq s \leq r_i \\ 1 \leq t \leq r_j}} d(V_i(s), V_j(t)) \leq d(X_{ij}^*, X_{ji}^*) + \frac{\varepsilon^5}{49}.$$

By the choice of parts of  $\mathcal{Q}$  within atoms, and because  $|V_i| \geq qm = 4^k m$ , we have

$$\begin{aligned} |X_{ij}^*| &\geq |X_{ij}| - 2^{k-1}m \\ &\geq |X_{ij}| \left(1 - \frac{2^k m}{\varepsilon |V_i|}\right) \\ &\geq |X_{ij}| \left(1 - \frac{1}{2^k \varepsilon}\right) \\ &\geq |X_{ij}| \left(1 - \frac{\varepsilon}{10}\right). \end{aligned}$$

So by the lemma, we know

$$|d(X_{ij}^*, X_{ji}^*) - d(X_{ij}, X_{ji})| < \frac{\varepsilon}{5}.$$

Recalling that

$$|d(X_{ij}, X_{ji}) - d(V_i, V_j)| > \varepsilon,$$

we have

$$\left| \frac{1}{r_i r_j} \sum_{\substack{1 \leq s \leq r_i \\ 1 \leq t \leq r_j}} d(V_i(s), V_j(t)) - d(V_i, V_j) \right| > \frac{3}{4}\varepsilon.$$

We can now allow our Cauchy–Schwarz inequality with  $n = q^2$  and  $m = r_i r_j$  gives

$$\begin{aligned} \frac{1}{q^2} \sum_{1 \leq s, t \leq q} d^2(V_i(s), V_j(t)) &\geq d^2(V_i, V_j) - \frac{\varepsilon^5}{25} + \frac{r_i r_j}{q^2} \cdot \frac{9\varepsilon^2}{16} \\ &\geq d^2(V_i, V_j) - \frac{\varepsilon^5}{25} + \frac{\varepsilon^4}{3}, \end{aligned}$$

using the fact that

$$\begin{aligned} \frac{r_i}{q} &\geq \left(1 - \frac{1}{m}\right) \frac{r_i}{q} \frac{m+1}{m} \geq \left(1 - \frac{1}{m}\right) \\ &\geq \frac{|X_{ij}^*|}{|V_i|} \geq \left(1 - \frac{1}{m}\right) \left(1 - \frac{\varepsilon}{10}\right) \frac{|X_{ij}|}{|V_i|} > \frac{4\varepsilon}{5}. \end{aligned}$$

Therefore

$$\begin{aligned} \text{ind}(\mathcal{Q}) &= \frac{1}{k^2 q^2} \sum_{\substack{1 \leq i < j \leq k \\ 1 \leq s, t \leq q}} d^2(V_i(s), V_j(t)) \\ &\geq \frac{1}{k^2} \sum_{1 \leq i < j \leq k} d^2(V_i, V_j) - \frac{\varepsilon^5}{25} + \varepsilon \binom{k}{2} \frac{\varepsilon^4}{3} \\ &\geq \text{ind}(P) + \frac{\varepsilon^5}{8}. \quad \square \end{aligned}$$

**Theorem.** Given an integer  $d$ , there exists  $c(d)$  such that

$$r(G) \leq c|G|$$

for every graph  $G$  with  $\Delta(G) \leq d$ .

*Proof.* Let  $t = R(d+1)$ . Pick  $\varepsilon \leq \min\left\{\frac{1}{t}, \frac{1}{2^d(d+1)}\right\}$ . Let  $\ell \geq t^2$ , and let  $L = L(\ell, \varepsilon)$ . We show that  $c = \frac{L}{\varepsilon}$  works.

Indeed, let  $G$  be a graph. Colour the edges of  $K_n$  by red and blue, where  $n \geq c|G|$ . Apply Szemerédi to the red graph with  $\ell, \varepsilon$  as above. Let  $H$  be the graph whose vertices are  $\{V_1, \dots, V_m\}$ , where  $V_1, \dots, V_m$  is the partition of the red graph. Let  $V_i, V_j \in E(H)$  if  $(V_i, V_j)$  is  $\varepsilon$ -uniform. Notice that  $m = |H| \geq \ell \geq t^2$ , and  $e(\bar{H}) \leq \varepsilon \binom{m}{2}$ . So  $H \supseteq K_t$ , or else by Turán's theorem, there are integers  $d_1, \dots, d_{t-1}$  and  $\sum d_i m$  and

$$e(\bar{H}) \geq \binom{d_i}{2} \geq (t-1) \binom{m/(t-1)}{2} \geq \varepsilon \binom{m}{2}$$

by our choice of  $\varepsilon$  and  $m$ .

We may as well assume all pairs  $V_i, V_j$  for  $1 \leq i < j \leq t$  are  $\varepsilon$ -uniform. We colour the edge of  $K_t$  green if  $d(V_i, V_j) \geq \frac{1}{2}$  (in the red graph), or white if  $< \frac{1}{2}$  (i.e. density  $> \frac{1}{2}$  in the blue graph).

By Ramsey's theorem, we may assume all pairs  $V_i V_j$  for  $1 \leq i < j \leq d+1$  are the same colour. We may wlog assume the colour is green, and we shall find a red  $G$  (a similar argument gives a blue  $G$  if the colour is white).

Indeed, take a vertex colouring of  $G$  with at most  $d+1$  colours (using  $\Delta(G) \leq d$ ), with no colour used more than  $|G|$  times. By the building lemma with  $H$  (in the lemma) being  $G$  (in this proof), and  $G$  (in the lemma) equal to the subgraph of the red graph spanned by  $V_1, \dots, V_{d+1}$  (here),

$$u = |V_i| \geq \frac{n}{L} \geq \frac{c|G|}{L} \geq \frac{|G|}{\varepsilon},$$

$r = d+1$ ,  $\lambda = \frac{1}{2}$ , and we are done. □

**Theorem** (Triangle removal lemma). Given  $\varepsilon > 0$ , there exists  $\delta > 0$  such that if  $|G| = n$  and  $G$  contains at most  $\delta n^3$  triangles, then there exists a set of at most  $\varepsilon n^2$  edges whose removal leaves no triangles.

*Proof.* Exercise. (See example sheet) □

**Corollary** (Roth, 1950's). Let  $\varepsilon > 0$ . Then if  $n$  is large enough, and  $A \subseteq [n] = \{1, 2, \dots, n\}$  with  $|A| \geq \varepsilon n$ , then  $A$  contains a 3-term arithmetic progression.

*Proof.* Define

$$B = \{(x, y) \in [2n]^2 : x - y \in A\}.$$

Then certainly  $|B| \geq \varepsilon n^2$ . We form a 3-partite graph with disjoint vertex classes  $X = [2n]$  and  $Y = [2n]$ ,  $Z = [4n]$ . If we have  $x \in X, y \in Y$  and  $z \in Z$ , we join  $x$  to  $y$  if  $(x, y) \in B$ ; join  $x$  to  $z$  if  $(x, z - x) \in B$  and join  $y$  to  $z$  if  $(z - y, y) \in B$ .

A triangle in  $G$  is a triple  $(x, y, z)$  where  $(x, y), (x, y + w), (x + w, y) \in B$  where  $w = z - x - y$ . Note that  $w < 0$  is okay. A 0-triangle is one with  $w = 0$ . There are at least  $\varepsilon n^2$  of these, one for each  $(x, y) \in B$ , and these are edge disjoint, because  $z = x + y$ .

Hence the triangles cannot be killed by removing  $\leq \varepsilon n^2/2$  edges. By the triangle removal lemma, we must have  $\geq \delta n^3$  triangles for some  $\delta$ . In particular, for  $n$  large enough, there is some triangle that is not a 0-triangle.

But then we are done, since

$$x - y - w, x - y, x - y + w \in A$$

where  $w \neq 0$ , and this is a 3-term arithmetic progression. □

## 5 Subcontraction, subdivision and linking

**Theorem** (Menger's theorem). Let  $G$  be a graph and  $s_1, \dots, s_k, t_1, \dots, t_k$  be distinct vertices. If  $\kappa(G) \geq k$ , then there exists  $k$  vertex disjoint paths from  $\{s_1, \dots, s_k\}$  to  $\{t_1, \dots, t_k\}$ .

**Lemma.** If  $\kappa(G) \geq 2k$  and  $G \supseteq TK_{5k}$ , then  $G$  is  $k$ -linked.

*Proof.* Let  $B$  be the set of "branch vertices" of  $TK_{5k}$ , i.e. the  $5k$  vertices joined by paths. By Menger's theorem, there exists  $2k$  vertex disjoint paths joining  $\{s_1, \dots, s_k, t_1, \dots, t_k\}$  to  $B$  (note that these sets might intersect). We say one of our  $2k$  paths *impinges* on a path in  $TK_{5k}$  if it meets that path and subsequently leaves it. Choose a set of  $2k$  paths impinging minimally (counting 1 per impingement).

Let these join  $\{s_1, \dots, t_k\}$  to  $\{v_1, \dots, v_{2k}\}$ , where  $B = \{v_1, \dots, v_{5k}\}$ . Then no path impinges on a path in  $TK_{5k}$  from  $\{v_1, \dots, v_{2k}\}$  to  $\{v_{2k+1}, \dots, v_{5k}\}$ . Otherwise, pick the path impinging closest to  $v_{2k+j}$ , and reroute it to  $v_{2k+j}$  rather than to something in  $\{v_1, \dots, v_{2k}\}$ , reducing the number of impingements.

Hence each path (of any  $2k$ ) meet at most one path in  $\{v_1, \dots, v_{2k}\}$  to  $\{v_{2k+1}, \dots, v_{5k}\}$  (once we hit it, we must stay there).

Thus, we may assume the paths from  $\{v_1, \dots, v_{2k}\}$  to  $\{v_{4k+1}, \dots, v_{5k}\}$  are not met at all. Then we are done, since we can link up  $v_j$  to  $v_{k+j}$  via  $v_{4k+j}$  to join  $s_j$  to  $t_k$ .  $\square$

**Lemma.** If  $e(G) \geq k|G|$ , then there exists some  $H$  with  $|H| \leq 2k$  and  $\delta(H) \geq k$  such that  $G \succ K_1 + H$ .

*Proof.* Consider the minimal subcontraction  $G'$  of  $G$  among those that satisfy  $e(G') \geq k|G'|$ . Then we must in fact have  $e(G') = k|G'|$ , or else we can just throw away an edge.

Since  $e(G') = k|G'|$ , it must be the case that  $\delta(G') \leq 2k$ . Let  $v$  be of minimum degree in  $G'$ . and set  $H = G'[\Gamma(v)]$ . Then  $G \succ K_1 + H$  and  $|H| \leq 2v$ .

To see that  $\delta(H) \geq k$ , suppose  $u \in V(H) = \Gamma(v)$ . Then by minimality of  $G'$ , we have

$$e(G'/uv) \leq k|G'| - k - 1.$$

But the number of edges we kill when performing this contraction is exactly 1 plus the number of triangles containing  $uv$ . So  $uv$  lies in at least  $k$  triangles of  $G'$ . In other words,  $\delta(H) \geq k$ .  $\square$

**Theorem.** If  $t \geq 3$  and  $e(G) \geq 2^{t-3}|G|$ , then  $G \succ K_t$ .

*Proof.* If  $t = 3$ , then  $G$  contains a cycle. So  $G \succ K_3$ . If  $t > 3$ , then  $G \succ K_1 + H$  where  $\delta(H) \geq 2^{t-3}$ . So  $e(H) \geq 2^{t-4}|H|$  and (by induction)  $H \succ K_{t-1}$ .  $\square$

**Lemma.** If  $\delta(G) \geq 2k$ , then  $G$  contains vertex disjoint subgraphs  $H, J$  with  $\delta(H) \geq k$ ,  $J$  connected, and every vertex in  $H$  has a neighbour in  $J$ .

*Proof.* We may assume  $G$  is connected, or else we can replace  $G$  by a component.

Now pick a subgraph  $J$  maximal such that  $J$  is connected and  $e(G/J) \geq k(|G| - |J| + 1)$ . Note that any vertex could be used for  $J$ . So such a  $J$  exist.

Let  $H$  be the subgraph spanned by the vertices having a neighbour in  $J$ . Note that if  $v \in V(H)$ , then  $v$  has at least  $k$  neighbours in  $H$ . Indeed, when we contract  $J \cup \{v\}$ , maximality tells us

$$e(G/(J \cup \{v\})) \leq k(|G| - |J| + 1) - k - 1,$$

and the number of edges we lose in the contraction is 1 plus the number of neighbours of  $v$  (it is easier to think of this as a two-step process — first contract  $J$ , so that we get an  $H + K_1$ , then contract  $v$  with the vertex coming from  $K_1$ ).  $\square$

**Theorem.** Let  $F$  be a graph with  $n$  edges and no isolated vertices. If  $\delta(G) \geq 2^n$ , then  $G \supseteq TF$ .

*Proof.* If  $n = 1$ , then this is immediate. If  $F$  consists of  $n$  isolated edges, then  $F \subseteq G$  (in fact,  $\delta(G) \geq 2n - 1$  is enough for this). Otherwise, pick an edge  $e$  which is not isolated. Then  $F - e$  has at most 1 isolated vertex. Apply the previous lemma to  $G$  to obtain  $H$  with  $\delta(H) \geq 2^{n-1}$ . Find a copy of  $T(F - e)$  in  $H$  (apart from the isolated vertex, if exists).

If  $e$  was just a leaf, then just add an edge (say to  $J$ ). If not, just construct a path going through  $J$  to act as  $e$ , since  $J$  is connected.  $\square$

**Lemma.** We have

$$c(t) \geq (\alpha + o(1))t\sqrt{\log t}.$$

To determine  $\alpha$ , let  $\lambda < 1$  be the root to  $1 - \lambda + 2\lambda \log \lambda = 0$ . In fact  $\lambda \approx 0.284$ . Then

$$\alpha = \frac{1 - \lambda}{2\sqrt{\log 1/\lambda}} \approx 0.319.$$

*Proof.* Consider a random graph  $G(n, p)$ , where  $p$  is a constant to be chosen later. So  $G(n, p)$  has  $n$  vertices and edges are chosen independently at random, with probability  $p$ . Here we fix a  $t$ , and then try to find the best combination of  $n$  and  $p$  to give the desired result.

A partition  $V_1, \dots, V_t$  of  $V(G)$  is said to be *complete* if there is an edge between  $V_i$  and  $V_j$  for all  $i$  and  $j$ . Note that having a complete partition is a *necessary*, but not sufficient condition for having a  $K_t$  minor. So it suffices to show that with probability  $> 0$ , there is no complete partition.

Writing  $q = 1 - p$ , a given partition with  $|V_i| = n_i$  is complete with probability

$$\begin{aligned} \prod_{i,j} (1 - q^{n_i n_j}) &\leq \exp\left(-\sum_{i < j} q^{n_i n_j}\right) \\ &\leq \exp\left(-\binom{t}{2} \prod q^{n_i n_j / \binom{t}{2}}\right) && \text{(AM-GM)} \\ &\leq \exp\left(-\binom{t}{2} q^{n^2/t^2}\right). \end{aligned}$$

The expected number of complete partitions is then

$$\leq t^n \exp\left(-\binom{t}{2} q^{n^2/t^2}\right),$$

As long as we restrict to the choices of  $n$  and  $q$  such that

$$t > n \sqrt{\frac{\log(1/q)}{\log n}},$$

we can bound this by

$$\leq \exp\left(n \log t - \binom{t}{2} \frac{1}{n}\right) = o(1)$$

in the limit  $t \rightarrow \infty$ . We set

$$q = \lambda, \quad n = \frac{t\sqrt{\log t}}{\sqrt{\log 1/\lambda}}.$$

Then with probability  $> 0$ , we can find a graph with only  $o(1)$  many complete partitions, and by removing  $o(n)$  many edges, we have a graph with

$$p\binom{n}{2} - o(n) = (\alpha + o(1))t\sqrt{\log t} \cdot n$$

many edges with no  $K_t$  minor.  $\square$

**Lemma.** Let  $k \in \mathbb{N}$  and  $G$  be a graph with  $e(G) \geq 11k|G|$ . Then there exists some  $H$  with

$$|H| \leq 11k + 2, \quad 2\delta(H) \geq |H| + 4k - 1$$

such that  $G \succ H$ .

*Proof.* We use our previous lemma as a starting point. Letting  $\ell = 11k$ , we know that we can find  $H_1$  such that

$$G \succ K_1 + H_1,$$

with  $|H_1| \leq 2\ell$  and  $\delta(H_1) \geq \ell$ . We shall improve the connectivity of this  $H_1$  at the cost of throwing away some elements.

By divine inspiration, consider the constant  $\beta \approx 0.37$  defined by the equation

$$1 = \beta \left(1 + \frac{\log 2}{\beta}\right),$$

and the function  $\phi$  defined by

$$\phi(F) = \beta\ell \frac{|F|}{2} \left(\log \frac{|F|}{\beta\ell} + 1\right).$$

Now consider the set of graphs

$$\mathcal{C} = \{F : |F| \geq \beta\ell, e(F) \geq \phi(F)\}.$$

Observe that  $H_1 \in \mathcal{C}$ , since  $\delta(H_1) \geq \ell$  and  $|H_1| \leq 2\ell$ .

Let  $H_2$  be a subcontraction of  $H_1$ , minimal with respect to  $\mathcal{C}$ .

Since the complete graph of order  $\lceil \beta\ell \rceil$  is not in  $\mathcal{C}$ , as  $\phi > \binom{\beta\ell}{2}$ , the only reason we are minimal is that we have hit the bound on the number of edges. Thus, we must have

$$|H_2| \geq \beta\ell + 1, \quad e(H_2) = \lceil \phi(H_2) \rceil,$$

and

$$e(H_2/uv) < \phi(H_2/uv)$$

for all edges  $uv$  of  $H_2$ .

Choose a vertex  $u \in H_2$  of minimum degree, and put  $H_3 = H_2[\Gamma(u)]$ . Then we have

$$|H_3| = \delta(H_2) \leq \frac{2\lceil \phi(H_2) \rceil}{|H_2|} \leq \left\lfloor \beta\ell \left( \log \left( \frac{|H_2|}{\beta\ell} \right) + 1 \right) + \frac{2}{|H_2|} \right\rfloor \leq \ell,$$

since  $\beta\ell \leq |H_2| \leq 2\ell$ .

Let's write  $b = \beta\ell$  and  $h = |H_2|$ . Then by the usual argument, we have

$$\begin{aligned} & 2\delta(H_3) - |H_3| \\ & \geq 2(\phi(H_2) - \phi(H_2/uv) - 1) - |H_3| \\ & \geq bh \left( \log \frac{h}{b} + 1 \right) - b(h-1) \left( \log \frac{h-1}{b} + 1 \right) - b \left( \log \frac{h}{b} + 1 \right) - 3 \\ & = b(h-1) \log \frac{h}{h-1} - 3 \\ & > b - 4, \end{aligned}$$

because  $h \geq b+1$  and  $x \log(1 + \frac{1}{x}) > 1 - \frac{1}{x}$  for real  $x > 1$ . So

$$2\delta(H_3) - |H_3| \geq \beta\ell - 4 > 4k - 4.$$

If we put  $H = K_2 + H_3$ , then  $G \succ H$ ,  $|H| \leq 11k + 2$  and

$$2\delta(H) - |H| \geq 2\delta(H_3) - |H_3| + 2 \geq 4k - 1. \quad \square$$

**Theorem.** We have

$$c(t) \leq 7t\sqrt{\log t}$$

if  $t$  is large.

*Proof.* For large  $t$ , we can choose  $k \in \mathbb{N}$  so that

$$5.8t\sqrt{\log_2 t} \leq 11k \leq 7t\sqrt{\log t}.$$

Let  $\ell = \lceil 1.01\sqrt{\log_2 t} \rceil$ . Then  $k \geq t\ell/2$ . We want to show that  $G \succ K_t$  if  $e(G) \geq 11k|G|$ .

We already know that  $G \succ H$  where  $h = |H| \leq 11k + 2$ , and  $2\delta(H) \geq |H| + 4k - 1$ . We show that  $H \succ K_t$ .

From now on, we work entirely in  $H$ . Note that any two adjacent vertices have  $\geq 4k+1$  common neighbours. Randomly select  $2t$  disjoint  $\ell$  sets  $U_1, \dots, U_{2t}$  in  $V(H)$ . This is possible since we need  $2t\ell$  vertices, and  $2t\ell \leq 4k$ .

Of course, we only need  $t$  many of those  $\ell$  sets, and we want to select the ones that have some chance of being useful to us.

Fix a part  $U$ . For any vertex  $v$  (not necessarily in  $U$ ), its degree is at least  $h/2$ . So the probability that  $v$  has no neighbour in  $U$  is at most  $\binom{h/2}{\ell} / \binom{h}{\ell} \leq 2^{-\ell}$ . Write  $X(U)$  for the vertices having no neighbour in  $U$ , then  $\mathbb{E}|X(U)| \leq 2^{-\ell}h$ .

We say  $U$  is *bad* if  $|X(U)| > 8 \cdot 2^{-\ell}h$ . Then by Markov's inequality, the probability that  $U$  is bad is  $< \frac{1}{8}$ . Hence the expected number of bad sets amongst



$U_1, \dots, U_{2t}$  is  $\leq \frac{t}{4}$ . By Markov again, the probability that there are more than  $\frac{t}{2}$  bad sets is  $< \frac{1}{2}$ .

We say a pair of sets  $(U, U')$  is *bad* if there is no  $U - U'$  edge. Now the probability that  $U, U'$  is bad is  $\mathbb{P}(U' \subseteq X(U))$ . If we condition on the event that  $U$  is good (i.e. not bad), then this probability is bounded by

$$\binom{8 \cdot 2^{-\ell} h}{\ell} / \binom{h - \ell}{\ell} \leq 8^\ell 2^{-\ell^2} \left(1 + \frac{\ell}{h - \ell}\right)^\ell \leq 8^\ell 2^{-\ell^2} e^{\ell^2/(h - \ell)} \leq 9^\ell 2^{-\ell^2}$$

if  $t$  is large (hence  $\ell$  is slightly large and  $h$  is very large). We can then bound this by  $\frac{1}{8t}$ .

Hence, the expected number of bad pairs (where one of them is good) is at most

$$\frac{1}{8t} \binom{2t}{2} \leq \frac{t}{4}.$$

So the probability that there are more than  $\frac{t}{2}$  such bad pairs is  $< \frac{1}{2}$ .

Now with positive probability,  $U_1, \dots, U_{2t}$  has at most  $\frac{t}{2}$  bad sets and at most  $\frac{t}{2}$  bad pairs amongst the good sets. Then we can find  $t$  many sets that are pairwise good. We may wlog assume they are  $U_1, \dots, U_t$ .

Fixing such a choice  $U_1, \dots, U_t$ , we now work deterministically. We would be done if each  $U_i$  were connected, but they are not in general. However, we can find  $W_1, \dots, W_t$  in  $V(H) \setminus (U_1 \cup \dots \cup U_t)$  such that  $U_i \cup W_i$  is connected for  $1 \leq i \leq t$ .

Indeed, if  $U_i = \{u_1, \dots, u_\ell\}$ , we pick a common neighbour of  $u_{j-1}$  and  $u_j$  if  $u_{j-1}u_j \notin E(H)$ , and put it in  $W_i$ . In total, this requires us to pick  $\leq t\ell$  distinct vertices in  $V(H) - (U_1 \cup \dots \cup U_t)$ , and we can do this because  $u_{j-1}$  and  $u_j$  have at least  $4k - t\ell \geq t\ell$  common neighbours in this set.  $\square$

**Lemma.** Let  $d \geq 0$ ,  $k \geq 2$  and  $h \geq d + \lfloor 3k/2 \rfloor$  be integers.

Let  $G$  be a graph,  $S = \{s_1, \dots, s_k\} \subseteq V(G)$ . Suppose there exists disjoint subgraphs  $C_1, \dots, C_h$  of  $G$  such that

- (\*) Each  $C_i$  is either connected, or each of its component meets  $S$ . Moreover, each  $C_i$  is adjacent to all but at most  $d$  of the  $C_j$ ,  $j \neq i$  not meeting  $S$ .
- (†) Moreover, no  $S$ -cut of order  $< k$  avoids  $d + 1$  of  $C_1, \dots, C_h$ .

Then  $G$  contains disjoint non-empty connected subgraphs  $D_1, \dots, D_m$ , where

$$m = h - \lfloor k/2 \rfloor,$$

such that for  $1 \leq i \leq k$ ,  $s_i \in D_i$ , and  $D_i$  is adjacent to all but at most  $d$  of  $D_{k+1}, \dots, D_m$ .

*Proof.* Suppose the theorem is not true. Then there is a minimal counterexample  $G$  (with minimality defined with respect to proper subgraphs).

We first show that we may assume  $G$  has no isolated vertices. If  $v$  is an isolated vertex, and  $v \notin S$ , then we can simply apply the result to  $G - v$ . If  $v \in S$ , then the  $S$ -cut  $(S, V(G) \setminus \{v\})$  of order  $k - 1$  avoids  $h - k \geq d - 1$  of the  $C_i$ 's.

**Claim.** If  $(A, B)$  is an  $S$ -cut of order  $k$  avoiding  $d + 1$  many of the  $C_i$ 's, then  $B = V(G)$  and  $A$  is discrete.

Indeed, given such an  $S$ -cut, we define

$$\begin{aligned} S' &= A \cap B \\ G' &= G[B] - E(S') \\ C'_i &= C_i \cap G'. \end{aligned}$$

We now make a further claim:

**Claim.**  $(G', S', \{C'_i\})$  satisfies the hypothesis of the lemma.

We shall assume this claim, and then come back to establishing the claim.

Assuming these claims, we show that  $G$  isn't a counterexample after all. Since  $(G', S', \{C'_i\})$  satisfies the hypothesis of the lemma, by minimality, we can find subgraphs  $D'_1, \dots, D'_m$  in  $G'$  that satisfy the conclusion of the lemma for  $G'$  and  $S'$ . Of course, these do not necessarily contain our  $s_i$ . Assuming these claims, we can construct the desired  $D_i$ . However, we can just add paths from  $S'$  to the  $s_i$ .

Indeed,  $G[A]$  has no  $S$ -cut  $(A'', B'')$  of order less than  $k$  with  $S' \subseteq B''$ , else  $(A'', B'' \cup B)$  is an  $S$ -cut of  $G$  of the forbidden kind, since  $A$ , and hence  $A''$  avoids too many of the  $C_i$ . Hence by Menger's theorem, there are vertex disjoint paths  $P_1, \dots, P_k$  in  $G[A]$  joining  $s_i$  to  $s'_i$  (for some labelling of  $S'$ ). Then simply take

$$D_i = D'_i \cup P_i.$$

To check that  $(G', S', \{C'_i\})$  satisfy the hypothesis of the lemma, we first need to show that the  $C'_i$  are non-empty.

If  $(A, B)$  avoids  $C_j$ , then by definition  $C_j \cap A = \emptyset$ . So  $C_j = C'_j$ . By assumption, there are  $d + 1$  many  $C_j$ 's for which this holds. Also, since these don't meet  $S$ , we know each  $C_i$  is adjacent to at least one of these  $C_j$ 's. So in particular,  $C'_i$  is non-empty since  $C'_i \supseteq C_i \cap C'_j = C_i \cap C_j$ .

Let's now check the conditions:

- (\*) For each  $C'_i$ , consider its components. If there is some component that does not meet  $S'$ , then this component is contained in  $B \setminus A$ . So this must also be a component of  $C_i$ . But  $C_i$  is connected. So this component is  $C_i$ . So  $C'_i = C_i$ . Thus, either  $C'_i$  is connected, or all components meet  $S'$ .

Moreover, since any  $C'_j$  not meeting  $S'$  equals  $C_j$ , it follows that each  $C_i$  and so each  $C'_i$  is adjacent to all but at most  $d$  of these  $C'_j$ 's. Therefore (\*) holds.

- (†) If  $(A', B')$  is an  $S'$ -cut in  $G'$  avoiding some  $C'_i$ , then  $(A \cup A', B')$  is an  $S$ -cut of  $G$  of the same order, and it avoids  $C_i$  (since  $C'_i$  doesn't meet  $S'$ , and we have argued this implies  $C'_i = C_i$ , and so  $C_i \cap A = \emptyset$ ). In particular, no  $S'$ -cut  $(A', B')$  of  $G'$  has order less than  $k$  and still avoids  $d + 1$  of  $C'_1, \dots, C'_h$ .

We can now use our original claim. We know what  $S$ -cuts of order  $k$  loop like, so we see that if there is an edge that does not join two  $C_i$ 's, then we can contract the edge to get a smaller counterexample. Recalling that  $G$  has no isolated vertices, we see that in fact  $V(G) = \cup V(C_i)$ , and  $|C_i| = 1$  unless  $C_i \subseteq S$ .

Let

$$C = \bigcup \{V(C_i) : |C_i| \geq 2\}.$$

We claim that there is a set  $I$  of  $|C|$  independent edges meeting  $C$ . If not, by Hall's theorem, we can find  $X \subseteq C$  whose neighbourhood  $Y$  (in  $V(G) - S$ ) such that  $|Y| < |X|$ . Then  $(S \cup Y, V(G) - X)$  is an  $S$ -cut of order  $|S| - |X| + |Y| < |S| = k$  avoiding  $\geq |G| - |S| - |Y|$  many  $C_i$ 's, and this is

$$\geq |G| - |S| - |X| + 1 \geq |G| - |C| - k + 1.$$

But  $h \leq |G| - |C| + \frac{|C|}{2}$ , since  $|G| - |C|$  is the number of  $C_i$ 's of size 1, so we can bound the above by

$$\geq h - \frac{|C|}{2} - k + 1 \geq h - \frac{3k}{2} + 1 \geq d + 1.$$

This contradiction shows that  $I$  exists.

Now we can just write down the  $D_i$ 's. For  $1 \leq i \leq k$ , set  $D_i$  to be  $s_i$  if  $s_i$  is a  $C_\ell$  with  $|C_\ell| = 1$ , i.e.  $s_i \notin C$ . If  $s_i \in C$ , let  $D_i$  be the edge of  $I$  meeting  $S_i$ .

Let  $D_{k+1}, \dots, D_m$  each be single vertices of  $G - S -$  (ends of  $I$ ). Note those exist because  $\geq |G| - k - |C| \geq h - \lfloor \frac{3k}{2} \rfloor = m - k$  vertices are available. Note that each  $D_i$  contains a  $C_\ell$  with  $|C_\ell| = 1$ . So each  $D_i$  is joined to all but  $\leq d$  many  $D_j$ 's.  $\square$

**Theorem.** Let  $G$  be a graph with  $\kappa(G) \geq 2k$  and  $e(G) \geq 11k|G|$ . Then  $G$  is  $k$ -linked. In particular,  $f(k) \leq 22k$ .

*Proof.* We know that under these conditions,  $G \succ H$  with  $|H| \leq 11k + 2$  and  $2\delta(H) \geq |H| + 4k - 1$ . Let  $h = |H|$ , and  $d = h - 1 - \delta(H)$ . Note that

$$h = 2d - 2 + 2\delta(H) - h \geq 2d + 4k.$$

Let  $C_1, \dots, C_k$  be connected subgraphs of  $G$  which contract to form  $H$ . Then clearly each  $C_i$  is joined to all but at most  $h - 1 - \delta(H) = d$  other  $C_j$ 's. Now let  $s_1, \dots, s_k, t_1, \dots, t_k$  be the vertices we want to link up. Let  $S = \{s_1, \dots, s_k, t_1, \dots, t_k\}$ . Observe that  $G$  has no  $S$ -cut of order  $< 2k$  at all since  $\kappa(G) \geq 2k$ . So the conditions are satisfied, but with  $2k$  instead of  $k$ .

So there are subgraphs  $D_1, \dots, D_m$ , where  $m = h - k \geq 2d + 3k$  as described. We may assume  $s_i \in D_i$  and  $t_i \in D_{k+i}$ . Note that for each pair  $(D_i, D_j)$ , we can find  $\geq m - 2d \geq 3k$  other  $D_\ell$  such that  $D_\ell$  is joined to both  $D_i$  and  $D_j$ . So for each  $i = 1, \dots, k$ , we can find an unused  $D_\ell$  not among  $D_1, \dots, D_{2k}$  that we can use to connect up  $D_i$  and  $D_{k+i}$ , hence  $s_i$  and  $t_i$ .  $\square$

**Lemma.** If  $\delta(a) \geq t^2$  and  $G$  is  $\binom{t+2}{3}$ -linked, then  $G \supseteq TK_t$ .

*Proof.* Since  $\delta(G) \geq t^2$ , we may pick vertices  $v_1, \dots, v_t$  and sets  $U_1, \dots, U_t$  all disjoint so that  $|U_i| = t - 1$  and  $v_i$  is joined to  $U_i$ . Then  $G - \{v_1, \dots, v_k\}$  is still  $\binom{t+2}{2} - t = \binom{t}{2}$ -linked. So  $U_1, \dots, U_{t-1}$  may be joined with paths to form a  $TK_t$  in  $G$ .  $\square$

**Lemma.** Let  $k, d \in \mathbb{N}$  with  $k \leq \frac{d+1}{2}$ , and suppose  $e(G) \geq d|G|$ . Then  $G$  contains a subgraph  $H$  with

$$e(H) = d|H| - kd + 1, \quad \delta(H) \geq d + 1, \quad \kappa(H) \geq k.$$

*Proof.* Define

$$\mathcal{E}_{d,k} = \{F \subseteq G : |F| \geq d, e(F) > d|F| - kd\}.$$

We observe that  $D \in \mathcal{E}_{d,k}$ , but  $K_d \notin \mathcal{E}_{d,k}$ . Hence  $|F| > d$  for all  $F \in \mathcal{E}_{d,k}$ . Let  $H$  be a subgraph of  $G$  minimal with respect to  $H \in \mathcal{E}_{d,k}$ . Then  $e(H) = d|H| - kd + 1$ , and  $\delta(H) \geq d + 1$ , else  $H - v \in \mathcal{E}_{d,k}$  for some  $v \in H$ .

We only have to worry about the connectivity condition. Suppose  $S$  is a cutset of  $H$ . Let  $C$  be a component of  $G - S$ . Since  $\delta(h) \geq d + 1$ , we have  $|C \cup S| \geq d + 2$  and  $|H - C| \geq d + 2$ .

Since neither  $C \cup S$  nor  $H - C$  lies in that class, we have

$$e(C \cup S) \leq d|C \cup S| - kd \leq d|C| + d|S| - kd$$

and

$$e(H - C) \leq d|H| - d|C| - kd.$$

So we find that

$$e(H) \leq d|H| + d|S| - 2kd.$$

But we know this is equal to  $d|H| - kd + 1$ . So we must have  $|S| \geq k$ .  $\square$

**Theorem.**

$$\frac{t^2}{16} \leq t(t) \leq 13 \binom{t+1}{2}.$$

*Proof.* The lower bound comes from (disjoint copies of)  $K_{t^2/8, t^2/8}$ . For the upper bound, write

$$d = 13 \binom{t+1}{2}, \quad k = t \cdot (t+1).$$

Then we can find a subgraph  $H$  with  $\delta(H) \geq d + 1$  and  $\kappa(H) \geq k + 1$ .

Certainly  $|H| \geq d + 2$ , and we know

$$e(H) \geq d|H| - kd > (d - k)|H| \geq 11 \binom{t+1}{2} |H|.$$

So we know  $H$  is  $\binom{t+1}{2}$ -linked, and so  $H \supseteq TK_t$ .  $\square$

## 6 Extremal hypergraphs

**Theorem** (Erdős). Let  $G$  be  $\ell$ -uniform of order  $n$  with  $p\binom{n}{\ell}$  edges, where  $p \geq 2n^{-1/t^{\ell-1}}$ . Then  $G$  contains  $K_{\ell}^{\ell}(t)$  (provided  $n$  is large).

*Proof.* Assume that  $t \geq 2$ . We proceed by induction on  $\ell \geq 2$ . For each  $(\ell-1)$ -set  $\sigma \subseteq V(G)$ , let

$$N(\sigma) = \{v : \sigma \cup \{v\} \in E(G)\}.$$

Then the average degree is

$$\binom{n}{\ell-1}^{-1} \sum |N(\sigma)| = \binom{n}{\ell-1}^{-1} \ell |E(G)| = p(n-\ell+1).$$

For each of the  $\binom{n}{t}$  many  $t$ -sets  $T \subseteq V(G)$ , let

$$D(T) = \{\sigma : T \subseteq N(\sigma)\}.$$

Then

$$\sum_T |D(T)| = \sum_{\sigma} \binom{|N(\sigma)|}{t} \geq \binom{n}{\ell-1} \binom{p(n-\ell+1)}{t},$$

where the inequality follows from convexity, since we can assume that  $p(n-\ell+1) \geq t-1$  if  $n$  is large.

In particular, there exists a  $T$  with

$$|D(T)| \geq \binom{n}{t}^{-1} \binom{n}{\ell-1} \binom{p(n-\ell+1)}{t} \geq \frac{1}{2} p^t \binom{n}{\ell-1}.$$

when  $n$  is large. By our choice of  $t$ , we can write this as

$$|D(T)| \geq 2^{t-1} n^{-1/t^{\ell-2}} \binom{n}{\ell-1}.$$

If  $\ell = 2$ , then this is  $\geq 2^{t-1} \geq t$ . So  $T$  is joined to a  $T$ -set giving  $K_{t,t}$ .

If  $\ell \geq 3$ , then  $|D(T)| \geq 2n^{-1/t^{\ell-2}} \binom{n}{\ell-1}$ . So, by induction, the  $(\ell-1)$ -uniform hypergraph induced by the  $\sigma$  with  $T \subseteq N(\sigma)$  contains  $K_{\ell-1}^{\ell-1}(t)$ , giving  $K_{\ell}^{\ell}(t)$  with  $T$ .  $\square$